

Systems with the integer rounding property in normal monomial subrings

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ABSTRACT

Let C be a clutter and let A be its incidence matrix. If the linear system $x \ge 0$; $xA \le 1$ has the integer rounding property, we give a description of the canonical module and the a-invariant of certain normal subrings associated to C. If the clutter is a connected graph, we describe when the aforementioned linear system has the integer rounding property in combinatorial and algebraic terms using graph theory and the theory of Rees algebras. As a consequence we show that the extended Rees algebra of the edge ideal of a bipartite graph is Gorenstein if and only if the graph is unmixed.

Key words: canonical module, *a*-invariant, normal ideal, perfect graph, maximal cliques, Rees algebra, Ehrhart ring, integer rounding property.

1 INTRODUCTION

A clutter \mathcal{C} with finite vertex set X is a family of subsets of X called edges, and none of which is included in another. The set of vertices and edges of \mathcal{C} are denoted by $V(\mathcal{C})$ and $E(\mathcal{C})$, respectively. A basic example of a clutter is a graph.

Let \mathcal{C} be a clutter with finite vertex set $X = \{x_1, \ldots, x_n\}$. We shall always assume that \mathcal{C} has no isolated vertices, i.e., each vertex occurs in at least one edge. Let f_1, \ldots, f_q be the edges of \mathcal{C} , and let $v_k = \sum_{x_i \in f_k} e_i$ be the *characteristic vector* of f_k , where e_i is the *ith* unit vector in \mathbb{R}^n . The *incidence matrix* A of \mathcal{C} is the $n \times q$ matrix with column vectors v_1, \ldots, v_q . If $a = (a_i)$ and $c = (c_i)$ are vectors in \mathbb{R}^n , then $a \le c$ means that $a_i \le c_i$ for all i. Thus, $a \ge 0$ means that $a_i \ge 0$ for all i. The system $x \ge 0$; $x A \le 1$ has the *integer rounding property* if

$$\lceil \min\{\langle y, \mathbf{1} \rangle | y \ge 0; \ Ay \ge \alpha\} \rceil = \min\{\langle y, \mathbf{1} \rangle | \ Ay \ge \alpha; \ y \in \mathbb{N}^q\}$$

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for each integral vector α for which $\min\{\langle y, \mathbf{1} \rangle | y \geq 0; Ay \geq \alpha\}$ is finite. Here **1** denotes the vector with all its entries equal to 1, and \langle , \rangle denotes the standard inner product. For a thorough study of this property see (Schrijver 1986).

Let $R = K[x_1, ..., x_n]$ be a polynomial ring over a field K, and let $w_1, ..., w_r$ be the set of all integral vectors α such that $0 \le \alpha \le v_i$ for some i. We will examine the integer rounding property using the monomial subring:

$$S = K[x^{w_1}t, \dots, x^{w_r}t] \subset R[t],$$

where t is a new variable. As usual, we use the notation $x^a := x_1^{a_1} \cdots x_n^{a_n}$, where $a = (a_1, \dots, a_n) \in \mathbb{N}^n$. The subring S is called *normal* if S is integrally closed, i.e., $S = \overline{S}$, where \overline{S} is the integral closure of S in its field of fractions; see (Vasconcelos 2005).

The contents of this paper are as follows. One of the results in (Brennan et al. 2008) shows that S is normal if and only if the system $x \ge 0$; $xA \le 1$ has the integer rounding property (see Theorem 2.1). As a consequence, we show that if all edges of C have the same number of elements and either linear system $x \ge 0$; $xA \le 1$ or $x \ge 0$; $xA \ge 1$ has the integer rounding property, then, the subring $K[x^{v_1}t, \ldots, x^{v_q}t]$ is normal (see Corollary 2.4).

Let G be a connected graph and let A be its incidence matrix. The main results of Section 3 show that the following conditions are equivalent (see Theorems 3.2 and 3.3):

- (a) $x \ge 0$; $xA \le 1$ has the integer rounding property.
- (b) $x \ge 0$; $x A \ge 1$ has the integer rounding property (see Definition 2.2).
- (c) $R[x^{v_1}t, \ldots, x^{v_q}t]$ is normal, where v_1, \ldots, v_q are the column vectors of A.
- (d) $K[x^{v_1}t, \ldots, x^{v_q}t]$ is normal.
- (e) $K[t, x_1t, \ldots, x_nt, x^{v_1}t, \ldots, x^{v_q}t]$ is normal.
- (f) The induced subgraph of the vertices of any two vertex disjoint odd cycles of G is connected.

The most interesting part of this result is the equivalence between (a), which is a linear optimization property, and (f), which is a graph theoretical property. Edge ideals are defined in Section 2. We prove that the ring in (e) is isomorphic to the extended Rees algebra of the edge ideal of G (see Proposition 3.1). If G is bipartite and I = I(G) is its edge ideal, we are able to prove that the extended Rees algebra of I is a Gorenstein standard K-algebra if and only if G is unmixed (see Corollary 4.3). If we work in the more general context of clutters, none of the conditions (a) to (e) are equivalent. Some of these conditions are equivalent under certain assumptions (Gitler et al. 2009).

In Section 4 we introduce the canonical module and the *a*-invariant of *S*. This invariant plays a key role in the theory of Hilbert functions (Bruns and Herzog 1997). There are some methods, based on combinatorial optimization, that have been used to study canonical modules of edge subrings of bipartite graphs (Valencia and Villarreal 2003). Our approach to study canonical modules is inspired by these

methods. If S is a normal domain, we express the canonical module of S and its a-invariant in terms of the vertices of the polytope

$$\{x | x \ge 0; xA \le 1\}$$

(see Theorem 4.1). We are able to give an explicit description of the canonical module of S and its a-invariant when C is the clutter of maximal cliques of a perfect graph (Theorem 4.2).

For unexplained terminology and notation on commutative algebra and integer programming, we refer to (Bruns and Herzog 1997, Vasconcelos 2005, Villarreal 2001) and (Schrijver 1986), respectively.

2 INTEGER ROUNDING AND NORMALITY

We continue using the definitions and terms from the Introduction. In what follows, \mathbb{N} denotes the set of non-negative integers and \mathbb{R}_+ denotes the set of non-negative real numbers. Let $\mathcal{A} \subset \mathbb{Z}^n$. The convex hull of \mathcal{A} is denoted by conv(\mathcal{A}), and the cone generated by \mathcal{A} is denoted by $\mathbb{R}_+\mathcal{A}$.

THEOREM 2.1 (Brennan et al. 2008). Let C be a clutter and let v_1, \ldots, v_q be the columns of the incidence matrix A of C. If w_1, \ldots, w_r is the set of all $\alpha \in \mathbb{N}^n$ such that $\alpha \leq v_i$ for some i, then the system $x \geq 0$; $xA \leq 1$ has the integer rounding property if and only if the subring $K[x^{w_1}t, \ldots, x^{w_r}t]$ is normal.

Next we give an application of this result, but first we need to introduce some more terminology and notation. We have already defined in the Introduction when the linear system $x \ge 0$; $xA \le 1$ has the integer rounding property. The following is a dual notion.

DEFINITION 2.2. Let A be a matrix with entries in \mathbb{N} . The system $x \ge 0$; $xA \ge 1$ has the *integer rounding* property if

$$\max\{\langle v, \mathbf{1} \rangle | Av < w; v \in \mathbb{N}^q\} = |\max\{\langle v, \mathbf{1} \rangle | v > 0; Av < w\}|$$

for each integral vector w for which the right hand side is finite.

Let $A = \{v_1, \dots, v_q\}$ be a set of points in \mathbb{N}^n , let P be the convex hull of A, and let $R = K[x_1, \dots, x_n]$ be a polynomial ring over a field K. The *Ehrhart ring* of the lattice polytope P is the monomial subring

$$A(P) = K[\{x^a t^i | a \in \mathbb{Z}^n \cap iP; i \in \mathbb{N}\}] \subset R[t],$$

where t is a new variable. A nice property of A(P) is that it is always a normal domain (Bruns and Herzog 1997). Let C be a clutter with vertex set $X = \{x_1, \ldots, x_n\}$ and edge set E(C). For use below recall that C is called *uniform* if all its edges have the same number of elements. The *edge ideal* of C, denoted by I(C), is the ideal of C generated by all monomials $\prod_{x_i \in C} x_i = x_e$ such that $e \in E(C)$. The *Rees algebra* of $C \in E(C)$, denoted by $C \in E(C)$, denoted by $C \in E(C)$.

$$R[It] := R \oplus It \oplus \cdots \oplus I^i t^i \oplus \cdots \subset R[t],$$

see (Vasconcelos 2005, Chapter 1) for a nice presentation of Rees algebras.

The next result holds for arbitrary monomial ideals (not necessarily square-free). This is the only place in the paper where a result is stated for arbitrary monomial ideals.

THEOREM 2.3 (Dupont and Villarreal 2010). Let $I = (x^{v_1}, \ldots, x^{v_q}) \subset R$ be a monomial ideal, and let A be the matrix with column vectors v_1, \ldots, v_q . Then, the system $x \ge 0$; $xA \ge 1$ has the integer rounding property if and only if R[It] is normal.

COROLLARY 2.4. Let C be a uniform clutter, let A be its incidence matrix, and let v_1, \ldots, v_q be the columns of A. If either system $x \ge 0$; $xA \le 1$ or $x \ge 0$; $xA \ge 1$ has the integer rounding property and $P = \text{conv}(v_1, \ldots, v_q)$, then

$$K[x^{v_1}t,\ldots,x^{v_q}t]=A(P).$$

PROOF. In general, the subring $K[x^{v_1}t, \ldots, x^{v_q}t]$ is contained in A(P). Assume that $x \geq 0$; $xA \leq 1$ has the integer rounding property, and that every edge of C has d elements. Let w_1, \ldots, w_r be the set of all $\alpha \in \mathbb{N}^n$ such that $\alpha \leq v_i$ for some i. Then by Theorem 2.1, the subring $K[x^{w_1}t, \ldots, x^{w_r}t]$ is normal. Using that v_1, \ldots, v_q is the set of w_i , with $|w_i| = d$, it is not hard to see that A(P) is contained in $K[x^{v_1}t, \ldots, x^{v_q}t]$.

Assume that $x \ge 0$; $xA \ge 1$ has the integer rounding property. Let $I = I(\mathcal{C})$ be the edge ideal of \mathcal{C} , and let R[It] be its Rees algebra. By Theorem 2.3, R[It] is a normal domain. Since the clutter \mathcal{C} is uniform, the required equality follows at once (Escobar et al. 2003, Theorem 3.15).

The converse of Corollary 2.4 fails as the following example shows.

EXAMPLE 2.5. Let C be the uniform clutter with vertex set $X = \{x_1, \dots, x_8\}$ and edge set

$$E(C) = \{\{x_3, x_4, x_6, x_8\}, \{x_2, x_5, x_6, x_7\}, \{x_1, x_4, x_5, x_8\}, \{x_1, x_2, x_3, x_8\}\}.$$

The characteristic vectors of the edges of C are

$$v_1 = (0, 0, 1, 1, 0, 1, 0, 1), v_2 = (0, 1, 0, 0, 1, 1, 1, 0),$$

 $v_3 = (1, 0, 0, 1, 1, 0, 0, 1), v_4 = (1, 1, 1, 0, 0, 0, 0, 1).$

Let A be the incidence matrix of C with column vectors v_1, \ldots, v_4 , and let P be the convex hull of $\{v_1, \ldots, v_4\}$. It is not hard to verify that the set

$$\{(v_1, 1), (v_2, 1), (v_3, 1), (v_4, 1)\}$$

is a Hilbert basis in the sense of (Schrijver 1986). Therefore, we have the equality

$$K[x^{v_1}t, x^{v_2}t, x^{v_3}t, x^{v_4}t] = A(P).$$

Using Theorem 2.3 and (Brennan et al. 2008, Theorem 2.12) it is seen that none of the two systems $x \ge 0$; $xA \le 1$ and $x \ge 0$; $xA \le 1$ have the integer rounding property.

3 INCIDENCE MATRICES OF GRAPHS

Let G be a connected graph with vertex set $X = \{x_1, \dots, x_n\}$, and let v_1, \dots, v_q be the column vectors of the incidence matrix A of G. The main result here is a purely combinatorial description of the integer

rounding property of the system $x \ge 0$; $xA \le 1$. Other equivalent algebraic conditions of this property will be presented.

Let $R = K[x_1, ..., x_n]$ be a polynomial ring over a field K, and let I = I(G) be the edge ideal of G. Recall that the *extended Rees algebra* of I is the subring

$$R[It, t^{-1}] := R[It][t^{-1}] \subset R[t, t^{-1}],$$

where R[It] is the Rees algebra of I. Rees algebras of edge ideals of graphs were first studied in (Simis et al. 1994).

PROPOSITION 3.1. $R[It, t^{-1}] \simeq K[t, x_1t, ..., x_nt, x^{v_1}t, ..., x^{v_q}t].$

PROOF. We set $S = K[t, x_1t, ..., x_nt, x^{v_1}t, ..., x^{v_q}t]$. Note that S and $R[It, t^{-1}]$ are both integral domains of the same Krull dimension; this follows from the dimension formula given in (Sturmfels 1996, Lemma 4.2). Thus, it suffices to prove that there is an epimorphism $\overline{\psi}: S \to R[It, I^{-1}]$ of K-algebras.

Let $u_0, u_1, \ldots, u_n, t_1, \ldots, t_q$ be a new set of variables, and let φ , ψ be the maps of K-algebras defined by the diagram

$$K[u_0, u_1, \dots, u_n, t_1, \dots, t_q] \xrightarrow{\psi} R[It, t^{-1}] \qquad u_0 \xrightarrow{\varphi} t, \qquad u_0 \xrightarrow{\psi} t^{-1},$$

$$\downarrow^{\varphi} \qquad \qquad \downarrow^{\psi} \qquad \qquad u_i \xrightarrow{\varphi} x_i t, \qquad u_i \xrightarrow{\psi} x_i,$$

$$\downarrow^{\varphi} \qquad \qquad \qquad t_i \xrightarrow{\varphi} x^{v_i} t, \qquad t_i \xrightarrow{\psi} x^{v_i} t.$$

To complete the proof, we will show that there is an epimorphism $\overline{\psi}$ of K-algebras that makes this diagram commutative, i.e., $\psi = \overline{\psi}\varphi$. To show the existence of $\overline{\psi}$, we need only to show the inclusion $\ker(\varphi) \subset \ker(\psi)$. As $\ker(\varphi)$ being a toric ideal is generated by binomials (Sturmfels 1996), it suffices to prove that any binomial of $\ker(\varphi)$ belongs to $\ker(\psi)$. Let

$$f = u_0^{a_0} u_1^{a_1} \cdots u_n^{a_n} t_1^{b_1} \cdots t_q^{b_q} - u_0^{c_0} u_1^{c_1} \cdots u_n^{c_n} t_1^{d_1} \cdots t_q^{d_q}$$

be a binomial in $ker(\varphi)$. Then,

$$t^{a_0}(x_1t)^{a_1}\cdots(x_nt)^{a_n}(x^{v_1}t)^{b_1}\cdots(x^{v_q}t)^{b_q}=t^{c_0}(x_1t)^{c_1}\cdots(x_nt)^{c_n}(x^{v_1}t)^{d_1}\cdots(x^{v_q}t)^{d_q}$$

Taking degrees in t and $x = \{x_1, \dots, x_n\}$, we obtain

$$a_0 + (a_1 + \dots + a_n) + (b_1 + \dots + b_q) = c_0 + (c_1 + \dots + c_n) + (d_1 + \dots + d_q),$$

$$a_1 + \dots + a_n + 2(b_1 + \dots + b_q) = c_1 + \dots + c_n + 2(d_1 + \dots + d_q).$$

Thus $-a_0 + b_1 + \cdots + b_q = -c_0 + d_1 + \cdots + d_q$, and we obtain the equality

$$t^{-a_0}x_1^{a_1}\cdots x_n^{a_n}(x^{v_1}t)^{b_1}\cdots (x^{v_q}t)^{b_q}=t^{-c_0}x_1^{c_1}\cdots x_n^{c_n}(x^{v_1}t)^{d_1}\cdots (x^{v_q}t)^{d_q},$$

i.e., $f \in ker(\psi)$, as required.

We come to the main result of this section.

THEOREM 3.2. Let G be a connected graph and let A be its incidence matrix. Then, the system

$$x \ge 0; xA \le 1$$

has the integer rounding property if and only if the induced subgraph of the vertices of any two vertex disjoint odd cycles of G is connected.

PROOF. Let v_1, \ldots, v_q be the column vectors of A. According to (Simis et al. 1998, Theorem 1.1, cf. Villarreal 2005, Corollary 3.10), the subring $K[Gt] := K[x^{v_1}t, \ldots, x^{v_q}t]$ is normal if and only if any two vertex disjoint odd cycles of G can be connected by at least one edge of G. Thus, we need only to show that K[Gt] is normal if and only the system $x \geq 0$; $xA \leq 1$ has the integer rounding property. Let I = I(G) be the edge ideal of G. Since G is connected, the subring K[Gt] is normal if and only if the Rees algebra R[It] of I is normal (Simis et al. 1998, Corollary 2.8). By a result of (Herzog et al. 1991), R[It] is normal if and only if $R[It, t^{-1}]$ is normal if and only if the subring

$$S = K[t, x_1t, \dots, x_nt, x^{v_1}t, \dots, x^{v_q}t]$$

is normal. Thus, we can apply Theorem 2.1 to conclude that S is normal if and only if the system $x \ge 0$; $xA \le 1$ has the integer rounding property.

THEOREM 3.3. Let G be a connected graph and let A be its incidence matrix. Then, the system $x \ge 0$; $xA \le 1$ has the integer rounding property if and only if any of the following equivalent conditions hold

- (a) x > 0; xA > 1 is a system with the integer rounding property.
- (b) R[It] is a normal domain, where I = I(G) is the edge ideal of G.
- (c) $K[x^{v_1}t, \ldots, x^{v_q}t]$ is normal, where v_1, \ldots, v_q are the column vectors of A.
- (d) $K[t, x_1t, \ldots, x_nt, x^{v_1}t, \ldots, x^{v_q}t]$ is normal.

PROOF. According to Theorem 2.3, the system $x \ge 0$; $xA \ge 1$ has the integer rounding property if and only if the Rees algebra R[It] is normal. Thus, the result follows from the proof of Theorem 3.2.

4 THE CANONICAL MODULE AND THE a-INVARIANT

In this section, we give a description of the canonical module and the *a*-invariant for subrings arising from systems with the integer rounding property.

Let C be a clutter with vertex set $X = \{x_1, \ldots, x_n\}$, and let v_1, \ldots, v_q be the columns of the incidence matrix A of C. For use below, consider the set w_1, \ldots, w_r of all $\alpha \in \mathbb{N}^n$ such that $\alpha \leq v_i$ for some i. Let $R = K[x_1, \ldots, x_n]$ be a polynomial ring over a field K and let

$$S = K[x^{w_1}t, \dots, x^{w_r}t] \subset R[t]$$

be the subring of R[t] generated by $x^{w_1}t, \ldots, x^{w_r}t$, where t is a new variable. As $(w_i, 1)$ lies in the hyperplane $x_{n+1} = 1$ for all i, S is a standard K-algebra. Thus, a monomial x^at^b in S has degree b. In what follows, we assume that S has this grading. Recall that the a-invariant of S, denoted a(S), is the degree as a rational function of the Hilbert series of S, see for instance (Villarreal 2001, p. 99). If S is Cohen-Macaulay and ω_S is the canonical module of S, then

$$a(S) = -\min\{i \mid (\omega_S)_i \neq 0\},\tag{1}$$

see (Bruns and Herzog 1997, p. 141) and (Villarreal 2001, Proposition 4.2.3). This formula applies if *S* is normal because normal monomial subrings are Cohen-Macaulay (Hochster 1972). If *S* is normal, then by a formula of Danilov-Stanley, see (Bruns and Herzog 1997, Theorem 6.3.5) and (Danilov 1978), the canonical module of *S* is the ideal given by

$$\omega_{\mathcal{S}} = (\{x^a t^b | (a, b) \in \mathbb{N}\mathcal{B} \cap (\mathbb{R}_+ \mathcal{B})^o\}), \tag{2}$$

where $\mathcal{B} = \{(w_1, 1), \dots, (w_r, 1)\}$ and $(\mathbb{R}_+\mathcal{B})^o$ is the interior of $\mathbb{R}_+\mathcal{B}$ relative to $aff(\mathbb{R}_+\mathcal{B})$, the affine hull of $\mathbb{R}_+\mathcal{B}$. In our case, $aff(\mathbb{R}_+\mathcal{B}) = \mathbb{R}^{n+1}$.

The next theorem complements a result of (Brennan et al. 2008). In loc. cit. a somewhat different expression for the canonical module and *a*-invariant are shown. Our expressions are simpler because they only involve the vertices of a certain polytope, while in (Brennan et al. 2008) some other parameters are involved.

THEOREM 4.1. Let C be a clutter with incidence matrix A, let v_1, \ldots, v_q be the columns of A, and let w_1, \ldots, w_r be the set of all $\alpha \in \mathbb{N}^n$ such that $0 \le \alpha \le v_i$ for some i. If the system $x \ge 0$; $xA \le 1$ has the integer rounding property and ℓ_1, \ldots, ℓ_m are the non-zero vertices of $P = \{x \mid x \ge 0; xA \le 1\}$, then the subring $S = K[x^{w_1}t, \ldots, x^{w_r}t]$ is normal, the canonical module of S is given by

$$\omega_{S} = \left(\left\{ x^{a} t^{b} \middle| (a, b) \begin{pmatrix} -\ell_{1} & \cdots & -\ell_{m} & e_{1} & \cdots & e_{n} \\ 1 & \cdots & 1 & 0 & \cdots & 0 \end{pmatrix} > 0 \right\} \right), \tag{3}$$

and the a-invariant of S is equal to $-(\max_i \{\lfloor |\ell_i| \rfloor\} + 1)$. Here $|\ell_i| = \langle \ell_i, 1 \rangle$.

PROOF. Note that in Eq. (3) we regard the ℓ_i 's and $e_i's$ as column vectors. The normality of S follows from Theorem 2.1. Let $P = \{x \mid x \ge 0; xA \le 1\}$ and let T(P) be its antiblocking polyhedron

$$T(P) := \{z \mid z \ge 0; \langle z, x \rangle \le 1 \text{ for all } x \in P\}.$$

By the finite basis theorem (Schrijver 1986), we can write

$$P = \{ z | z > 0; \langle z, w_i \rangle < 1 \,\forall i \} = \text{conv}(\ell_0, \ell_1, \dots, \ell_m), \tag{4}$$

where $\ell_1, \ldots, \ell_m, \ell_0$ are the vertices of P and $\ell_0 = 0$. Notice that the vertices of P are in \mathbb{Q}^n_+ . From Eq. (4), we readily get the equality

$$\{z \mid z \ge 0; \langle z, \ell_i \rangle \le 1 \,\forall i\} = T(P). \tag{5}$$

Using Eq. (4) again and noticing that $\langle \ell_i, w_i \rangle \leq 1$ for all i, j, we get

$$\mathbb{R}^n_+ \cap (\operatorname{conv}(\ell_0, \dots, \ell_m) + \mathbb{R}_+ \{-e_1, \dots, -e_n\}) = \{z \mid z \ge 0; \langle z, w_i \rangle \le 1 \,\forall i\}.$$

Hence, using this equality and (Schrijver 1986, Theorem 9.4), we obtain

$$\mathbb{R}_{+}^{n} \cap (\text{conv}(w_{1}, \dots, w_{r}) + \mathbb{R}_{+}\{-e_{1}, \dots, -e_{n}\}) = \{z \mid z \geq 0; \langle z, \ell_{i} \rangle \leq 1 \,\forall i\}. \tag{6}$$

By (Fulkerson 1971, Theorem 8), we have the equality

$$conv(w_1, ..., w_r) = \mathbb{R}^n_+ \cap (conv(w_1, ..., w_r) + \mathbb{R}_+ \{-e_1, ..., -e_n\}).$$

Therefore, using Eqs. (5) and (6), we conclude the following duality:

$$P = \{x \mid x \ge 0; \langle x, w_i \rangle \le 1 \,\forall i\} = \operatorname{conv}(\ell_0, \ell_1, \dots, \ell_m),$$

$$\operatorname{conv}(w_1, \dots, w_r) = \{x \mid x \ge 0; \langle x, \ell_i \rangle \le 1 \,\forall i\} = T(P).$$
(7)

We set $\mathcal{B} = \{(w_1, 1), \dots, (w_r, 1)\}$. Note that $\mathbb{Z}\mathcal{B} = \mathbb{Z}^{n+1}$. From Eq. (7), it is seen that

$$\mathbb{R}_{+}\mathcal{B} = H_{e_{1}}^{+} \cap \dots \cap H_{e_{n}}^{+} \cap H_{(-\ell_{1},1)}^{+} \cap \dots \cap H_{(-\ell_{n},1)}^{+}. \tag{8}$$

Here H_a^+ denotes the closed halfspace $H_a^+ = \{x \mid \langle x, a \rangle \geq 0\}$ and H_a stands for the hyperplane through the origin with normal vector a. Notice that

$$H_{e_1} \cap \mathbb{R}_+ \mathcal{B}, \ldots, H_{e_n} \cap \mathbb{R}_+ \mathcal{B}, H_{(-\ell_1, 1)} \cap \mathbb{R}_+ \mathcal{B}, \ldots, H_{(-\ell_m, 1)} \cap \mathbb{R}_+ \mathcal{B}$$

are proper faces of $\mathbb{R}_+\mathcal{B}$. Hence, from Eq. (8), we get that a vector (a, b), with $a \in \mathbb{R}^n$, $b \in \mathbb{R}$, is in the relative interior of $\mathbb{R}_+\mathcal{B}$ if and only if the entries of a are positive and $\langle (a, b), (-\ell_i, 1) \rangle > 0$ for all i. Thus, the required expression for ω_S , i.e., Eq. (3), follows using the normality of S and the Danilov-Stanley formula given in Eq. (2).

It remains to prove the formula for a(S), the a-invariant of S. Consider the vector $(1, b_0)$, where $b_0 = \max_i \{\lfloor |\ell_i| \rfloor\} + 1$. Using Eq. (3), it is not hard to see (by direct substitution of $(1, b_0)$) that the monomial $x^1t^{b_0}$ is in ω_S . Thus, from Eq. (1), we get $a(S) \geq -b_0$. Conversely, if the monomial x^at^b is in ω_S , then again from Eq. (3) we get $\langle (-\ell_i, 1), (a, b) \rangle > 0$ for all i and $a_i \geq 1$ for all i, where $a = (a_i)$. Hence,

$$b > \langle a, \ell_i \rangle \ge \langle \mathbf{1}, \ell_i \rangle = |\ell_i| \ge \lfloor |\ell_i| \rfloor.$$

Since b is an integer, we obtain $b \ge \lfloor |\ell_i| \rfloor + 1$ for all i. Therefore, $b \ge b_0$, i.e., $\deg(x^a t^b) = b \ge b_0$. As $x^a t^b$ was an arbitrary monomial in ω_s , by the formula for the a-invariant of S given in Eq. (1), we obtain that $a(S) \le -b_0$. Altogether one has $a(S) = -b_0$, as required.

MONOMIAL SUBRINGS OF CLIQUES OF PERFECT GRAPHS

Let S be a set of vertices of a graph G, the *induced subgraph* $\langle S \rangle$ is the maximal subgraph of G with vertex set S. A *clique* of a graph G is a subset of the set of vertices that induces a complete subgraph. Let G be a graph with vertex set $X = \{x_1, \ldots, x_n\}$. A *colouring* of the vertices of G is an assignment of colours to the vertices of G in such a way that adjacent vertices have distinct colours. The *chromatic number* of G is the minimal number of colours in a colouring of G. A graph is *perfect* if, for every induced subgraph G, the chromatic number of G equals the size of the largest complete subgraph of G. Let G be a subset of the vertices of G. The set G is called *independent* if no two vertices of G are adjacent.

For use below we consider the empty set as a clique whose vertex set is empty. The *support* of a monomial x^a is given by $\sup(x^a) = \{x_i \mid a_i > 0\}$. Note that $\sup(x^a) = \emptyset$ if and only if a = 0.

THEOREM 4.2. Let G be a perfect graph and let $S = K[x^{\omega_1}t, \dots, x^{\omega_r}t]$ be the subring generated by all square-free monomials x^at such that $supp(x^a)$ is a clique of G. Then, the canonical module of S is given by

$$\omega_{S} = \left(\left\{ x^{a}t^{b} \middle| (a,b) \left(\begin{array}{cccc} -a_{1} & \cdots & -a_{m} & e_{1} & \cdots & e_{n} \\ 1 & \cdots & 1 & 0 & \cdots & 0 \end{array} \right) \geq \mathbf{1} \right\} \right),$$

where a_1, \ldots, a_m are the characteristic vectors of the maximal independent sets of G, and the a-invariant of S is equal to $-(\max_i \{|a_i|\} + 1)$.

PROOF. Let v_1, \ldots, v_q be the set of characteristic vectors of the maximal cliques of G. Note that w_1, \ldots, w_r is the set of all $\alpha \in \mathbb{N}^n$ such that $\alpha \leq v_i$ for some i. Since G is a perfect graph, by (Korte and Vygen 2000, Theorem 16.14) we have the equality

$$P = \{x | x \ge 0; x A \le 1\} = \text{conv}(a_0, a_1, \dots, a_p),$$

where $a_0 = 0$ and a_1, \ldots, a_p are the characteristic vectors of the independent sets of G. We may assume that a_1, \ldots, a_m correspond to the maximal independent sets of G. Furthermore, since P has only integral vertices, by a result of (Lovász 1972), the system $x \ge 0$; $xA \le 1$ is totally dual integral, i.e., the minimum in the LP-duality equation

$$\max\{\langle \alpha, x \rangle | x \ge 0; xA \le 1\} = \min\{\langle y, 1 \rangle | y \ge 0; Ay \ge \alpha\}$$
(9)

has an integral optimum solution y for each integral vector α with finite minimum. In particular the system $x \ge 0$; $xA \le 1$ satisfies the integer rounding property. Therefore, the result follows readily from Theorem 4.1.

For use below recall that a graph G is called *unmixed* if all maximal independent sets of G have the same cardinality. Unmixed bipartite graphs have been nicely characterized in (Villarreal 2007).

COROLLARY 4.3. Let G be a connected bipartite graph, and let I = I(G) be its edge ideal. Then, the extended Rees algebra $R[It, t^{-1}]$ is a Gorenstein standard K-algebra if and only if G is unmixed.

PROOF. Let ω_S be the canonical module of $S = R[It, t^{-1}]$. Recall that S is Gorenstein if and only if ω_S is a principal ideal (Bruns and Herzog 1997). Since any bipartite graph is a perfect graph, the result follows using Proposition 3.1 together with the description of the canonical module given in Theorem 4.2.

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RESUMO

Seja \mathcal{C} uma desordem (família de Sperner) e seja A sua matriz de incidência. Se o sistema linear $x \geq 0$; $xA \leq 1$ tem a propriedade do arredondamento inteiro, fornecemos a descrição do módulo canônico e do a-invariante de certos subaneis monomiais associados a \mathcal{C} . Se a desordem é um grafo conexo, descreve-se quando o supra-mencionado sistema linear tem a propriedade do arredondamento inteiro em termos combinatórios e algébricos, usando a teoria dos grafos e a teoria das álgebras de Rees. Como consequência, mostra-se que a álgebra de Rees estendida do ideal de arestas de um grafo bipartido é um anel de Gorenstein se e somente se o grafo é de altura pura.

Palavras-chave: módulo canônico, *a*-invariante, ideal normal, grafo perfeito, cliques maximais, álgebra de Rees, anel de Erhart, propriedade do arredondamento inteiro.

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